## Example of Gomory's Cutting Plane Method

Consider the linear program

$$
\min \quad 2 x_{1}+15 x_{2}+18 x_{3}
$$

Subject to

$$
\begin{gathered}
-x_{1}+2 x_{2}-6 x_{3} \leq-10 \\
x_{2}+2 x_{3} \leq 6 \\
2 x_{1}+10 x_{3} \leq 19 \\
-x_{1}+x_{2} \leq-2 \\
x_{1}, x_{2}, x_{3} \geq 0
\end{gathered}
$$

We can solve this problem the dual simplex method algorithm. The final tableau is as follows:

|  | $x_{1}$ | $x_{2}$ | $x_{3}$ | $z_{1}$ | $z_{2}$ | $z_{3}$ | $z_{4}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $x_{1}$ | 1 | 10 | 0 | 5 | 0 | 3 | 0 | 7 |
| $z_{2}$ | 0 | 5 | 0 | 2 | 1 | 1 | 0 | 5 |
| $x_{3}$ | 0 | -2 | 1 | -1 | 0 | $-1 / 2$ | 0 | $1 / 2$ |
| $z_{4}$ | 0 | 11 | 0 | 5 | 0 | 3 | 1 | 5 |
|  | 0 | 31 | 0 | 8 | 0 | 3 | 0 | -23 |

Which corresponds to the solution

$$
\begin{aligned}
& x_{1}=7 \\
& x_{2}=0 \\
& x_{3}=\frac{1}{2}
\end{aligned} \quad \text { Objective function }=23
$$

Now, imagine that we need $x_{1}, x_{2}$ and $x_{3}$ to be integers. The third row reads

$$
\begin{gathered}
-2 x_{2}+x_{3}-z_{1}-\frac{1}{2} z_{3}=\frac{1}{2} \\
-2 x_{2}+x_{3}-z_{1}-\left\lfloor\frac{1}{2}\right\rfloor z_{3} \leq \frac{1}{2} \\
-2 x_{2}+x_{3}-z_{1}-z_{3} \leq \frac{1}{2}
\end{gathered}
$$

If all the variables are integers, we also have that

$$
\begin{aligned}
& -2 x_{2}+x_{3}-z_{1}-z_{3} \leq\left\lfloor\frac{1}{2}\right\rfloor \\
& -2 x_{2}+x_{3}-z_{1}-z_{3} \leq 0
\end{aligned}
$$

We can add this new inequality to our tableau in the form

$$
-2 x_{2}+x_{3}-z_{1}-z_{3}+z_{5}=0
$$

Inserting this inequality into our tableau, we obtain

|  | $x_{1}$ | $x_{2}$ | $x_{3}$ | $z_{1}$ | $z_{2}$ | $z_{3}$ | $z_{4}$ | $z_{5}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $x_{1}$ | 1 | 10 | 0 | 5 | 0 | 3 | 0 | 0 | 7 |
| $z_{2}$ | 0 | 5 | 0 | 2 | 1 | 1 | 0 | 0 | 5 |
| $x_{3}$ | 0 | -2 | 1 | -1 | 0 | $-1 / 2$ | 0 | 0 | $1 / 2$ |
| $z_{4}$ | 0 | 11 | 0 | 5 | 0 | 3 | 1 | 0 | 5 |
| $z_{5}$ | 0 | -2 | 1 | -1 | 0 | -1 | 0 | 1 | 0 |
|  | 0 | 31 | 0 | 8 | 0 | 3 | 0 | 0 | -23 |

The matrix in the body of the tableau is $A_{B}^{-1} A$. Thus, those columns corresponding to basis variables should give $A_{B}^{-1} A_{B}=I$. Looking at the highlighted basis columns above, this is clearly not the case in the $x_{3}$ column.

We fix this by subtracting the $x_{3}$ row from the last row:

|  | $x_{1}$ | $x_{2}$ | $x_{3}$ | $z_{1}$ | $z_{2}$ | $z_{3}$ | $z_{4}$ | $z_{5}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $x_{1}$ | 1 | 10 | 0 | 5 | 0 | 3 | 0 | 0 | 7 |
| $z_{2}$ | 0 | 5 | 0 | 2 | 1 | 1 | 0 | 0 | 5 |
| $x_{3}$ | 0 | -2 | 1 | -1 | 0 | $-1 / 2$ | 0 | 0 | $1 / 2$ |
| $z_{4}$ | 0 | 11 | 0 | 5 | 0 | 3 | 1 | 0 | 5 |
| $z_{5}$ | 0 | 0 | 0 | 0 | 0 | $-1 / 2$ | 0 | 1 | $-1 / 2$ |
|  | 0 | 31 | 0 | 8 | 0 | 3 | 0 | 0 | -23 |

$z_{5}$ is negative, and we therefore pivot on that row. The only column with a negative entry in that row is $z_{3}$, so we pivot there.

|  | $x_{1}$ | $x_{2}$ | $x_{3}$ | $z_{1}$ | $z_{2}$ | $z_{3}$ | $z_{4}$ | $z_{5}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $x_{1}$ | 1 | 10 | 0 | 5 | 0 | 0 | 0 | 6 | 4 |
| $z_{2}$ | 0 | 5 | 0 | 2 | 1 | 0 | 0 | 2 | 4 |
| $x_{3}$ | 0 | -2 | 1 | -1 | 0 | 0 | 0 | -1 | 1 |
| $z_{4}$ | 0 | 11 | 0 | 5 | 0 | 0 | 1 | 6 | 2 |
| $z_{5}$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | -2 | 1 |
|  | 0 | 31 | 0 | 8 | 0 | 0 | 0 | 6 | -26 |

This solution is both primal and dual optimal, with only integer solution.

$$
\begin{aligned}
& x_{1}=4 \\
& x_{2}=0 \\
& x_{3}=1
\end{aligned} \quad \text { Objective function }=26
$$

